

# Introduction to openMP and Shared-Memory Programming

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# What is shared-memory?

- Parallel programs have multiple execution units talking to each other
- Shared-memory has this communication happening in the memory of the computer
- This means that all of the execution units need to be on one physical machine (with some exceptions)
- Forms:
  - openMP
  - POSIX threads
  - Home made memory sharing

# How does OpenMP work?

- Compiler directives in the source code
- If the compiler doesn't recognize them, they act as comments
- In C/C++ these are pragmas  
`#pragma omp parallel`
- In FORTRAN these are directives  
`!$OMP PARALLEL`

# How do we compile?

- Compiling simply requires a command line option

|               |             |
|---------------|-------------|
| PGI Compilers | -mp         |
| Sun Studio 12 | -openmp -O3 |
| GNU 4         | -fopenmp    |

# How do we run?

- We need to tell the program how many instances to run
  - Bash – `export OMP_NUM_THREADS=4`
  - Csh – `setenv OMP_NUM_THREADS 4`
- Is OpenMP allowed to change this?
  - Bash – `export OMP_DYNAMIC=FALSE`
  - Csh – `setenv OMP_DYNAMIC TRUE`
- Can we nest the parallel parts?
  - Bash – `export OMP_NESTED=TRUE`
  - Csh – `setenv OMP_NESTED FALSE`

# Learning about ourselves

- We may need to know how many processors and how many threads there are  
`int omp_get_num_procs()`  
`int omp_get_num_threads()`
- Find out who we are  
`int omp_get_thread_num()`
- Set the number of threads to use  
`void omp_set_num_threads(int t)`
- FORTRAN has the same subroutines

# Loops – for (C/C++)

- Most common form of parallelism is through the use of loops
  - #pragma omp parallel for
  - for (i = first; i < size; i += prime) marked[i] = 1;
- At runtime, we need to know how many iterations
- This means we can't use break, return or exit from inside the loop, or goto's that lead to outside of the loop
- Continue is OK since the loop count doesn't change

# Loops – do (FORTRAN)

- In FORTRAN, the most common loop is the do loop

```
!$omp do  
    do loop statements  
!$omp end do
```

# Privatizing Variables

- Normally, all variables are shared among threads

```
#pragma omp parallel for
  for (i=0; i<max; i++)
    for (j=0; j<max2; j++)
      do_something
```

- The variable j will get badly abused by access from all threads

# Privatizing Variables

- We can make variables private to fix this

```
#pragma omp parallel for private(j)
  for (i=0; i<max; i++)
    for (j=0; j<max2; j++)
      do_something
```

- A new variable j will get created for each thread

# Privatizing Variables

- What if we want the value to enter the loop?  
#pragma omp parallel for private(j) firstprivate(j)  
for (i=0; i<max; i++)  
for (j=0; j<max2; j++)  
do\_something
- What if we want the value to leave the loop?  
#pragma omp parallel for private(j) lastprivate(j)

# Critical Sections

- What if we have a statement in our loop that needs to be done sequentially  
    `area += add_on;`
- In this statement, `area` is read and added to. If this is done by two threads at the same time, they read in the same value, then add their “`add_on`” value.
- We miss one of the “`add_on`” values – race condition
- How do we fix this?

# Critical Sections

```
#pragma omp parallel for private(x)
  for (i=0; i<n; i++) {
    x = (i+0.5)/n;
    #pragma omp critical
      area += 4.0/(1.0 + x*x);
  }
```

- This change puts a lock on the line calculating area so that only one thread has access at a time

```
!$omp critical
```

```
...
```

```
!$omp end critical
```

# General parallelism

- What if we need to have some chunk of code executed by each thread?

```
#pragma omp parallel
{
    do_subroutine();
    a = b + c;
}
```

- We can use private here as well to get private variables within the block
- We can use critical here if we have code that is sensitive to a race condition

# How to drop out of parallel code

- What if we have a statement in a parallel section that should only be executed by one thread  
e.g. a print statement should only be printed once

```
#pragma omp parallel
{
    ...
    #pragma omp single
    printf("The answer is %d", answer);
    ...
}
```

# Coordination

- Sometimes threads need to meet at a common point so that they can coordinate
- We can set a barrier – threads can't get past the barrier until everyone gets there

!\$omp barrier

#pragma omp barrier

# POSIX Threads

- Threads are more low level then OpenMP
  - The programmer is responsible for more work
  - This means the programmer also has much more control
- 
- `pthread_create`
  - `pthread_join`
  - `pthread_exit`
  - `pthread_mutex_lock`
  - `pthread_mutex_unlock`

# Creating a thread

- Threads need to be created and given the code to execute

```
int pthread_create(pthread_t *tid, const pthread_attr_t  
    *attr, void * (*func)(void *), void *arg);
```

- The first argument holds the resulting thread ID
- The second argument holds the thread attributes
- The third argument is a pointer to the function
- The fourth argument is a pointer to an argument

# Working inside the thread

- Who are we?  
`pthread_t pthread_self()`
- Once we're done, how do we finish?  
`void pthread_exit(void *status);`
- The pointer to status needs to point to a global variable, since all local variables disappear when the thread exits

# Waiting for a thread to finish

- Once we spawn off a thread, we may want to wait for it to finish

```
int pthread_join(pthread_t tid, void **status);
```

- This blocks until the thread finishes, so don't use it unless you really need it

# Critical sections again

- We can get race conditions with POSIX threads when we try and read/write from global variables
- It's the programmer's responsibility to deal with it

```
int pthread_mutex_lock(pthread_mutex_t *mptr);
int pthread_mutex_unlock(pthread_mutex_t *mptr);
```
- If you need to initialize/reinitialize a mutex

```
pthread_mutex_t mptr =
    PTHREAD_MUTEX_INITIALIZER
int pthread_mutex_init(pthread_mutex_t *mptr,
    pthread_mutexattr_t *attr):
```